Semantic models of higher order probability theory

Sam Staton, Oxford

partly jointly with Hongseok Yang, Ohad Kammar, Chris Heunen, Frank Wood...

Motivation

What is a general notion of statistical model?

How can we understand the fundamental structure?

Aims of semantic models:

- 1. Fundamental;
- 2. Towards applications in Machine Learning.

Bayes' law $P(x \mid d) = \frac{P(d \mid x) \times P(x)}{P(d)}$

Bayes' law

P(x | d) =



Bayes' law

$P(x \mid d) \propto P(d \mid x) \times P(x)$ Posterior \propto Likelihood \times Prior

Idealized Anglican

 $P(x \mid d) \propto P(d \mid x) \times P(x)$ Posterior \propto Likelihood \times Prior

Idealized Anglican = sequential programming + normalize observe sample



Idealized Anglican

 $P(x \mid d) \propto P(d \mid x) \times P(x)$ Posterior \propto Likelihood \times Prior

Idealized Anglican = sequential programming + normalize observe sample

> other languages: Church,Venture... Hakaru ... Stan...

http://www.robots.ox.ac.uk/~fwood/anglican/

Example

- 1. A call centre operator doesn't know what day it is.
- 2. He knows: weekends: avg 3 calls per hour. weekdays: avg 10 calls per hour.
- 3. He notices a 15 minute gap between calls.
- 4. Is it the weekend?

Example

- 1. A call centre operator doesn't know what day it is.
- 2. He knows: weekends: avg 3 calls per hour. weekdays: avg 10 calls per hour.
- 3. He notices a 15 minute gap between calls.
- 4. Is it the weekend?

```
normalize(
    let weekend = sample(bernoulli(2/7)) in
    let rate = if weekend then 3 else 10 in
    observe 0.25 from exp-dist(rate);
    return(weekend) )
```

Example

- 1. A call centre operator doesn't know what day it is.
- 2. He knows: weekends: avg 3 calls per hour. weekdays: avg 10 calls per hour.
- 3. He notices a 15 minute gap between calls.



Motivation

 $P(x \mid d) \propto P(d \mid x) \times P(x)$ Posterior \propto Likelihood \times Prior

Idealized Anglican = sequential programming + normalize observe sample

Aims of semantic models:

- 1. Justify program transformations;
- 2. Understand complex probabilistic models through semantics.

Main theorem: Commutativity

$$\begin{bmatrix} \operatorname{let} x = t \operatorname{in} \\ | \operatorname{let} y = u \operatorname{in} \\ | v \end{bmatrix} = \begin{bmatrix} \operatorname{let} y = u \operatorname{in} \\ | \operatorname{let} x = t \operatorname{in} \\ | v \end{bmatrix}$$

$$\underbrace{v}_{where x \text{ not free in } u, where x \text{ not free in } u$$

Aims of semantics:

y not free in t

- 1. Justify program transformations;
- 2. Understand complex probabilistic models through semantics.

Main theorem: Commutativity



where x not free in u, y not free in t



connections with multicategories, monoidal categories etc..

Overview of the rest of part 1

Compositional probability theory; a compositional theory of impropriety

- Summary of semantics
- Examples
- s-finite kernels and proof of commutativity.

To begin: a first order typed language.

Types:distributions over A $\mathbb{A}, \mathbb{B} ::= \mathbb{R} \mid \mathsf{P}(\mathbb{A}) \mid 1 \mid \mathbb{A} \times \mathbb{B} \mid \sum_{i \in I} \mathbb{A}_i$ countable*i* finite products*i* finite products

Terms: sequencing: let x = t in u normalize observe sample

For now: a first order typed language without recursion.

Types interpreted as measurable spaces.

Closed terms of type A interpreted as measures on [A].

Types interpreted as measurable spaces.

A measurable space (X, Σ_X) is a set *X* together with set Σ_X of subsets closed under countable unions and complements.

Types interpreted as measurable spaces.

e.g. $[real] = (\mathbb{R}, \Sigma_{\mathbb{R}})$

where $\Sigma_{\mathbb{R}}$ = Borel sets

generated by intervals e.g. (-2.2, 4.1) & countable unions & complements

A measurable space (X, Σ_X) is a set *X* together with set Σ_X of subsets closed under countable unions and complements.

Types interpreted as measurable spaces.

e.g. $[real] = (\mathbb{R}, \Sigma_{\mathbb{R}})$



where $\Sigma_{\mathbb{R}}$ = Borel sets

generated by intervals e.g. (-2.2, 4.1) & countable unions & complements

A measurable space (X, Σ_X) is a set *X* together with set Σ_X of subsets closed under countable unions and complements.

Types interpreted as measurable spaces.

e.g. $[real] = (\mathbb{R}, \Sigma_{\mathbb{R}})$

Closed terms interpreted as measures

Open terms $\Gamma \vdash t \colon \mathbb{A}$ are interpreted as kernels $\llbracket t \rrbracket \colon \llbracket \Gamma \rrbracket \times \Sigma_{\llbracket \mathbb{A} \rrbracket} \to [0, \infty].$ *measurable in* $\llbracket \Gamma \rrbracket$ $\exists t \rrbracket \colon \llbracket \Gamma \rrbracket \to \Sigma_{\llbracket \mathbb{A} \rrbracket} \to [0, \infty].$

Types interpreted as measurable spaces.

e.g. $[real] = (\mathbb{R}, \Sigma_{\mathbb{R}})$

Closed terms interpreted as measures

Open terms interpreted as kernels

Sequencing is integration

$$\llbracket \operatorname{\mathsf{let}} x = t \operatorname{\mathsf{in}} u \rrbracket(\gamma, U) \stackrel{\text{\tiny def}}{=} \int \llbracket u \rrbracket(\gamma, x, U) \llbracket t \rrbracket(\gamma, \mathrm{d} x)$$

Types interpreted as measurable spaces.

e.g. real type interpreted as \mathbb{R} with Borel sets

Closed terms interpreted as measures

Open terms interpreted as kernels

Sequencing is integration

"
$$[[\det x = t \text{ in } u]] = \int [[u]] d[[t]]]$$
"

see Kozen, 1981 Börgstrom, Gordon et al., ESOP 2011 Shan, Ramsey, POPL 2017

Closed terms interpreted as measures Sequencing is integration

"
$$\llbracket \operatorname{let} x = t \operatorname{in} u \rrbracket = \int \llbracket u \rrbracket \operatorname{d} \llbracket t \rrbracket$$
"

Key Theorem: Commutativity.

 $\begin{bmatrix} \operatorname{let} x = t \text{ in} \\ \operatorname{let} y = u \text{ in} \\ v \end{bmatrix} = \begin{bmatrix} \operatorname{let} y = u \text{ in} \\ \operatorname{let} x = t \text{ in} \\ v \end{bmatrix}$ Probabilistic programs \approx s-finite kernels. $\begin{pmatrix} \left(\int \llbracket v \rrbracket d\llbracket u \rrbracket \right) d\llbracket t \rrbracket \end{bmatrix} = \int \left(\int \llbracket v \rrbracket d\llbracket t \rrbracket \right) d\llbracket u \rrbracket$

Overview of the rest of part 1

Compositional probability theory; a compositional theory of impropriety

- Summary of semantics
- Examples
- s-finite kernels and proof of commutativity.

- 1. A call centre operator doesn't know what day it is.
- 2. He knows: weekends: avg 3 calls per hour. weekdays: avg 10 calls per hour.
- 3. He notices 4 calls in a given hour.
- 4. Is it the weekend?



20

Unnormalized posterior:

 $m(weekend=true) = 2/7 \times 0.168 = 0.048$

 $m(weekend=false) = 5/7 \times 0.019 = 0.014$



Unnormalized posterior:

 $m(weekend=true) = 2/7 \times 0.168 = 0.048$

 $m(weekend=false) = 5/7 \times 0.019 = 0.014$














```
normalize(
let weekend = sample(bernoulli(2/7)) in
let rate = if weekend then 3 else 10 in
observe 4 from poisson(rate);
return(weekend) )
```



normalize(let weekend = sample(bernoulli(2/7)) in let rate = if weekend then 3 else 10 in observe 4 from poisson(rate); return(weekend))

false



$$\left[\operatorname{\mathsf{let}} x = t \text{ in } u\right] = \int_{\left[\!\left[\mathbb{A} \right]\!\right]} \left[\!\left[u \right]\!\right] \, \mathrm{d}\left[\!\left[t \right]\!\right]$$

false

normalize(

let weekend = sample(bernoulli(2/7)) in let rate = if weekend then 3 else 10 in observe 4 from poisson(rate); return(weekend))



observe 4 from poisson(rate); return(weekend))



normalize(

let weekend = sample(bernoulli(2/7)) in let rate = if weekend then 3 else 10 in observe 4 from poisson(rate); return(weekend))

- 1. A call centre operator doesn't know what day it is.
- 2. He knows: weekends: avg 3 calls per hour. weekdays: avg 10 calls per hour.
- 3. He notices a 15 minute gap between calls.
- 4. Is it the weekend? exp(10)3 $\mathbf{2}$ 1 exp(3)0 normalize($15 \mathrm{mins}$ 0 60 mins let weekend = sample(bernoulli(2/7)) in let rate = if weekend then 3 else 10 in observe 0.25 from exp-dist(rate); return(weekend))

Unnormalized posterior:

 $m(weekend=true) = 2/7 \times 1.42 = 0.405$

 $m(weekend=false) = 5/7 \times 0.82 = 0.586$



- 1. A call centre operator doesn't know what time it is.
- 2. He knows how the avg num of calls varies with time.
- 3. He notices a 15 minute gap between calls.
- 4. What time is it?

```
normalize(
let time = sample(uniform(0,24)) in
let rate = f(time) in
observe 0.25 from exp-dist(rate);
return(weekend) )
```

Improper posteriors



let x = sample(Normal(0,1)) in let r = 1/f(x) in observe d from exp-dist(r); return x



Overview of the rest of part 1

Compositional probability theory; a compositional theory of impropriety

- Summary of semantics
- Examples
- s-finite kernels and proof of commutativity.

Overview of the rest of part 1

Compositional probability theory; a compositional theory of impropriety

- Summary of semantics
- Examples
- s-finite kernels and proof of commutativity.

Defn. A kernel is a function $\llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ such that for all $U \in \Sigma_{\llbracket A \rrbracket}, \llbracket t \rrbracket (-, U) : \llbracket \Gamma \rrbracket \to [0, \infty]$ is measurable and $\llbracket t \rrbracket (\gamma, -) : \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ is additive.

Defn. A kernel is a function $\llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ such that for all $U \in \Sigma_{\llbracket A \rrbracket}, \llbracket t \rrbracket (-, U) : \llbracket \Gamma \rrbracket \to [0, \infty]$ is measurable and ...

Defn. A kernel is a function $\llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ such that for all $U \in \Sigma_{\llbracket A \rrbracket}, \llbracket t \rrbracket (-, U) : \llbracket \Gamma \rrbracket \to [0, \infty]$ is measurable and ...

Sequencing is integration: $\llbracket \text{let } x = t \text{ in } u \rrbracket(\gamma, U) \stackrel{\text{\tiny def}}{=} \int \llbracket u \rrbracket(\gamma, x, U) \llbracket t \rrbracket(\gamma, dx)$

Defn. A kernel is a function $\llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ such that for all $U \in \Sigma_{\llbracket A \rrbracket}, \llbracket t \rrbracket (-, U) : \llbracket \Gamma \rrbracket \to [0, \infty]$ is measurable and ...

Sequencing is integration: $\llbracket \text{let } x = t \text{ in } u \rrbracket(\gamma, U) \stackrel{\text{def}}{=} \int \llbracket u \rrbracket(\gamma, x, U) \llbracket t \rrbracket(\gamma, dx)$

Is this measurable in γ ?

Defn. A kernel is a function $\llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ such that for all $U \in \Sigma_{\llbracket A \rrbracket}, \llbracket t \rrbracket (-, U) : \llbracket \Gamma \rrbracket \to [0, \infty]$ is measurable and ...

Sequencing is integration: $\llbracket \text{let } x = t \text{ in } u \rrbracket(\gamma, U) \stackrel{\text{def}}{=} \int \llbracket u \rrbracket(\gamma, x, U) \llbracket t \rrbracket(\gamma, dx)$

Is this measurable in γ ?

Answer: **unknown** (if t and u are allowed to range over arbitrary kernels.)

Commutativity Terms are interpreted as kernels.

 $\begin{bmatrix} \text{let } x = t \text{ in} \\ \text{let } y = u \text{ in} \\ v \end{bmatrix} = \begin{bmatrix} \text{let } y = u \text{ in} \\ \text{let } x = t \text{ in} \\ v \end{bmatrix}$ reordering lines

- Useful for program transformations.
- Essence of probability theory.

Commutativity Terms are interpreted as kernels. $\begin{bmatrix} \text{let } x = t \text{ in} \\ \text{let } y = u \text{ in} \\ v \end{bmatrix} = \begin{bmatrix} \text{let } y = u \text{ in} \\ \text{let } x = t \text{ in} \\ v \end{bmatrix}$ reordering lines Roughly amounts to $\int \left(\int \llbracket v \rrbracket d\llbracket u \rrbracket \right) d\llbracket t \rrbracket = \int \left(\int \llbracket v \rrbracket d\llbracket t \rrbracket \right) d\llbracket u \rrbracket$

Fubini's theorem

Commutativity Terms are interpreted as kernels. $\begin{vmatrix} \text{let } x = t \text{ in} \\ \text{let } y = u \text{ in} \\ v \end{vmatrix} = \begin{vmatrix} \text{let } y = u \text{ in} \\ \text{let } x = t \text{ in} \\ v \end{vmatrix}$ reordering lines Roughly amounts to $\int \left(\int \llbracket v \rrbracket d\llbracket u \rrbracket \right) d\llbracket t \rrbracket = \int \left(\int \llbracket v \rrbracket d\llbracket t \rrbracket \right) d\llbracket u \rrbracket$ Fubini's theorem

Warning: Fubini's theorem does not hold for arbitrary kernels.

Commutativity Terms are interpreted as kernels. $\begin{bmatrix} \text{let } x = t \text{ in} \\ \text{let } y = u \text{ in} \\ v \end{bmatrix} = \begin{bmatrix} \text{let } y = u \text{ in} \\ \text{let } x = t \text{ in} \\ v \end{bmatrix}$ reordering lines Roughly amounts to $\int \left(\int \llbracket v \rrbracket d\llbracket u \rrbracket \right) d\llbracket t \rrbracket = \int \left(\int \llbracket v \rrbracket d\llbracket t \rrbracket \right) d\llbracket u \rrbracket$

Fubini's theorem

Warning: Fubini's theorem does not hold for arbitrary kernels.

CommutativityTerms are interpreted as s-finite
kernels.
$$\begin{bmatrix} let x = t in \\ let y = u in \\ v \end{bmatrix} = \begin{bmatrix} let y = u in \\ let x = t in \\ v \end{bmatrix}$$
Roughly amounts toreordering
lines
$$\int (\int [v]] d[u]) d[t] = \int (\int [v]] d[t]) d[u]$$
Fubini's theorem

Warning: Fubini's theorem does not hold for arbitrary kernels.

s-finite kernels

Definition. A kernel $k : \llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ is **bounded** if $\exists n. \forall \gamma. \forall U. k(\gamma, U) < n.$

Definition. A kernel $k : \llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ is **s-finite** if it is a countable sum of bounded kernels.

$$\int \left(\int \llbracket v \rrbracket d\llbracket u \rrbracket \right) d\llbracket t \rrbracket = \int \left(\int \llbracket v \rrbracket d\llbracket t \rrbracket \right) d\llbracket u \rrbracket$$

Fubini's theorem

Warning: Fubini's theorem does not hold for arbitrary kernels.

s-finite kernels

Definition. A kernel $k : \llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ is **bounded** if $\exists n. \forall \gamma. \forall U. k(\gamma, U) < n.$

Definition. A kernel $k : \llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ is **s-finite** if it is a countable sum of bounded kernels.

Warning: Fubini's theorem does not hold for arbitrary kernels.

s-finite kernels

Definition. A kernel $k : \llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ is **bounded** if $\exists n. \forall \gamma. \forall U. k(\gamma, U) < n.$

Definition. A kernel $k : \llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ is **s-finite** if it is a countable sum of bounded kernels.

Theorem. A kernel is s-finite if and only if it is definable.

Warning: Fubini's theorem does not hold for arbitrary kernels.

Commutativity

Theorem.

$$\begin{bmatrix} \operatorname{let} x = \operatorname{t} & \operatorname{in} \\ \operatorname{let} y = \operatorname{u} & \operatorname{in} \\ v & \end{bmatrix} = \begin{bmatrix} \operatorname{let} y = \operatorname{u} & \operatorname{in} \\ \operatorname{let} x = \operatorname{t} & \operatorname{in} \\ v & \end{bmatrix} \\ \int \left(\int \begin{bmatrix} A_1 \\ \bullet \end{bmatrix} d \begin{bmatrix} u \end{bmatrix} e d \begin{bmatrix} t \end{bmatrix} = \int \left(\int \begin{bmatrix} v \end{bmatrix} d \begin{bmatrix} t \end{bmatrix} \right) d \begin{bmatrix} u \end{bmatrix}$$



Defn. A kernel is a function $\llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ such that for all $U \in \Sigma_{\llbracket A \rrbracket}, \llbracket t \rrbracket (-, U) : \llbracket \Gamma \rrbracket \to [0, \infty]$ is measurable and ...

Sequencing is integration: $\llbracket \text{let } x = t \text{ in } u \rrbracket(\gamma, U) \stackrel{\text{def}}{=} \int \llbracket u \rrbracket(\gamma, x, U) \llbracket t \rrbracket(\gamma, dx)$

Is this measurable in γ ?

Defn. A kernel is a function $\llbracket \Gamma \rrbracket \times \Sigma_{\llbracket A \rrbracket} \to [0, \infty]$ such that for all $U \in \Sigma_{\llbracket A \rrbracket}, \llbracket t \rrbracket (-, U) : \llbracket \Gamma \rrbracket \to [0, \infty]$ is measurable and ...

Sequencing is integration: $\llbracket \text{let } x = t \text{ in } u \rrbracket(\gamma, U) \stackrel{\text{def}}{=} \int \llbracket u \rrbracket(\gamma, x, U) \llbracket t \rrbracket(\gamma, dx)$

Is this measurable in γ ?

Answer: **yes** (if t and u are probabilistic programs, they are s-finite kernels.)

Summary

Theorem.

$$\begin{bmatrix} \text{let } x = t \text{ in} \\ \text{let } y = u \text{ in} \\ v \end{bmatrix} = \begin{bmatrix} \text{let } \frac{A_1}{Y} = t \text{ in} \\ \text{let } x = t \text{ in} \\ v \end{bmatrix}$$





- Probabilistic programs ¬^t have a measure-theoretic semantics.
- Commutativity holds.
- Probabilistic programs are s-finite kernels.
Overview

- **Part 1:** Compositional probability theory; a compositional theory of impropriety
- Part 2: Higher order functions; random functions as measures on a space of functions.

Overview

Summary of semantics

- Examples
- s-finite kernels and proof of commutativity.

Overview of part 2

Higher order functions; random functions as measures on a space of functions.

- Examples of regression and higher order functions
- Quasi-Borel spaces

Regression











normalize(let s = sample (normal 0 2) b = sample (normal 0 6) f = λx. s x + b in return f)



Samples from the prior

normalize(let s = sample (normal 0 2) 10 b = sample (normal 0 6) $f = \lambda x_{\bullet} s x + b in$ 9 observe 0.6 from (normal (f 0) .5) 8 observe 0.7 from (normal (f 1) .5) observe 1.2 from (normal (f 2) .5) 7 – observe 3.2 from (normal (f 3) .5) 6 observe 6.8 from (normal (f 4) .5) 5 observe 8.2 from (normal (f 5) .5) observe 8.4 from (normal (f 6) .5) 4 return **f**) 3 -2 -1 -2.5 6.0 0.0 0.5 1.0 1.5 2.0 3.0 3.5 4.0 4.5 5.0 5.5 Samples from the posterior



```
normalize(
let f =
  (let s = sample (normal 0 2)
       b = sample (normal 0 6) in
       return \lambda x. s x + b) in
observe 0.6 from (normal (f 0) .5)
observe 0.7 from (normal (f 1) .5)
observe 1.2 from (normal (f 2) .5)
observe 3.2 from (normal (f 3) .5)
observe 6.8 from (normal (f 4) .5)
observe 8.2 from (normal (f 5) .5)
observe 8.4 from (normal (f 6) .5)
return f)
```

More higher-order functions

```
normalize(
let f = add-change-points
  (let s = sample (normal 0 2)
       b = sample (normal 0 6) in
       return \lambda x_{\bullet} s x + b in
observe 0.6 from (normal (f 0) .5)
observe 0.7 from (normal (f 1) .5)
observe 1.2 from (normal (f 2) .5)
observe 3.2 from (normal (f 3) .5)
observe 6.8 from (normal (f 4) .5)
observe 8.2 from (normal (f 5) .5)
observe 8.4 from (normal (f 6) .5)
return f )
```

Posterior



Technical problem

Measure theory doesn't support HO fns well.

$$\operatorname{ev}: (\mathbb{R} \to_{\mathrm{m}} \mathbb{R}) \times \mathbb{R} \to \mathbb{R}, \quad \operatorname{ev}(f, x) = f(x).$$

[Aumann 61] ev is not measurable no matter which σ -algebra is used for $\mathbb{R} \rightarrow_m \mathbb{R}$.

[Corollary] The category of measurable spaces is not cartesian closed.

Overview of part 2 Higher order functions; random functions as measures on a space of functions.

- Examples of regression and higher order functions
- Quasi-Borel spaces

What about higher-order functions?

 $\mathbb{R} \rightarrow \mathbb{R}$ this is **not** a measurable space [Aumann 1961]

Easy to deal with operationally. But denotationally?

Staton, Yang, Heunen, Kammar, Wood, LICS 2016 Börgstrom, Dal Lago, Gordon, Szymczak, ICFP 2016

What about function types?



Models first order language with sample, Score Models higher order language with sample, Score

Theorem. Adequacy.

What about function types?



Models first order language with sample, SCOR Models higher order language with sample, SCORE

Theorem. Adequacy.

What about function types?



Models first order language with sample, score

Problem: $\mathbb{R} \to \mathbb{R}$ doesn't exist i.e. Hom($-\times \mathbb{R}$, \mathbb{R}): Meas^{op} \to Set isn't representable

What about function types ?** Ammar, Wood, LICS 2016

StandardYonedaSheaves onBorel spacesembeddingmeasurable spaces

Models first order language with sample, SCOR

What about function types ?**

StandardYonedaSheaves onBorel spacesembeddingmeasurable spaces

Models first order language with sample, score

(Functors $Meas^{op} \rightarrow Set$ that preserve countable products)

What about function types ?** Ammar, Wood, LICS 2016



What about function types ?* also Power, TCS 2016



"Random elements first."

Random elements a in X

$\alpha:\Omega\to X$

- X set of values.
- Ω set of random seeds.
- Random seed generator.

Random elements a in X in classical measure theory

$\alpha: \Omega \to X$

• X - set of values.

is a random element if $\alpha^{-1}(A) \in \Sigma$ for all $A \in \Theta$

- Ω set of random seeds.
- Random seed generator.

Random elements a in X in quasi-Borel spaces

$\alpha: \Omega \to X$

- X set of values.
- $\Omega = \mathbb{R}$ set of random seeds.
- Random seed generator.

Random elements a in X in quasi-Borel spaces

$\label{eq:alpha} \begin{array}{l} \alpha:\Omega \to X \\ & \mbox{is a random variable} \end{array}$ • X - set of values. if $\alpha {\in} M$

- $\Omega = \mathbb{R}$ set of random seeds.
- Random seed generator.

1. \mathbb{R} as random source

2. Borel subsets $\mathfrak{B} \subseteq 2^{\mathbb{R}}$

3. M ⊆ [ℝ→X]

Defn. A quasi-Borel space is a pair (X,M) where

- X is a set
- $M \subseteq [\mathbb{R} \rightarrow X]$ s.t.

Defn. A quasi-Borel space is a pair (X,M) where

- X is a set
- $M \subseteq [\mathbb{R} \to X]$

such that

- if $f : \mathbb{R} \to \mathbb{R}$ measurable and $g \in M$ then $gf \in M$.
- piecewise combination: if $\mathbb{R}= \bigcup_{i \in \mathbb{N}} R_i$ with R_i Borel and $\alpha_1, \alpha_2, \ldots \in M$, then $\bigcup_{i \in \mathbb{N}} (\alpha_i \cap (R_i \times X)) \in M$.
- all constant functions are in M

A morphism $(X, M) \rightarrow (Y, N)$ is a function $f : X \rightarrow Y$ such that $g \in M$ implies $fg \in N$

Heunen, Kammar,

Defn. A quasi-Borel space is a pair (X,M) where

- X is a set
- $M \subseteq [\mathbb{R} \rightarrow X]$ s.t. ...

Example: X is a standard Borel measurable space, $M \subseteq [\mathbb{R} \rightarrow X]$ comprises the measurable functions. Then 'morphism' = 'measurable function'.

A morphism $(X, M) \rightarrow (Y, N)$ is a function $f : X \rightarrow Y$ such that $g \in M$ implies $fg \in N$

Heunen, Kammar, Staton,

Defn. A quasi-Borel space is a pair (X,M) where

- X is a set
- $M \subseteq [\mathbb{R} \rightarrow X]$ s.t. ...

Example: X is a standard Borel measurable space, $M \subseteq [\mathbb{R} \rightarrow X]$ comprises the measurable functions.

Proposition. The category of quasi-Borel spaces is cartesian closed with countable sums.

Corollary. If a term t has first order type then [[t]] is a measurable function.

Heunen, Kammar, Staton,
Heunen, Kammar, Staton, Yang, LICS 2017





What is a measure on $X \rightarrow Y$?



Defn. A quasi-Borel space is a pair (X,M) where

- *X* is a set
- $M \subseteq [\mathbb{R} \rightarrow X]$ s.t. ...

Defn. A **measure** on a quasi-Borel space is a pair (μ, f) a fⁿ $f : \mathbb{R} \to X$ in M

(modulo inducing the same integration operator)



Proposition. A measure on $[X \rightarrow Y]$ is a pair



– a 'random function'.

normalize(let s = sample (normal 0 2) b = sample (normal 0 6) g = λx. s x + b in return g)



normalize(let s = sample (normal 0 2) b = sample (normal 0 6) g = λx. s x + b in return g)





Quasi-Borel spaces work well:

- Simple theorems e.g. "randomization lemmas" can be stated in the internal logic;
- A version of de Finetti's theorem for exchangeable sequences holds;
- One can justify higher order inference algorithms, (Metropolis-Hastings, Sequential Monte Carlo) jww Scibior, Vakar, Kai, Ostermann, Moss, Ghahramani

Motivation

What is a general notion of statistical model?

How can we understand the fundamental structure?

Aims of semantic models:

- 1. Fundamental;
- 2. Towards applications in Machine Learning.

What's next:

We're now looking at "exchangeable random structures" towards a general theory of Aldous-Hoover results and the connection with abstract types.

• Discussions with Ackerman, Freer, Roy, Bloom-Reddy.

Aims of semantic models:

- 1. Fundamental;
- 2. Towards applications in Machine Learning.

Overview

- **Part 1:** Compositional probability theory; a compositional theory of impropriety
- **Part 2:** Higher order functions; random functions as measures on a space of functions.